

Complete Parsimony Haplotype Inference Problem and Algorithms

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Haplotypes and genotypes

- **Haplotypes:** set of nucleotides in physical proximity on a chromosome strand.
- **Genotypes:** confluents of **haplotype pairs**.

HIPP

- Identifying individual haplotypes in a laboratory setting: feasible only for small studies.
- **Observation:** Number of unique haplotypes in a given population is small.
- ↪ **Minimize** number of **unique haplotypes**.
- ↪ **Haplotype Inference by Pure Parsimony (HIPP)**

Haplotypes and genotypes

- Let $m \in \mathbb{N}$.
- **Haplotype:** $h \in \{0, 1\}^m$
- **Genotype:** $g \in \{0, 1, 2\}^m$
- **Explanation** of genotypes by haplotypes:

h, h' explain g , if for each $i = 1, 2, \dots, m$:

- 1 $g_i = 0 \Rightarrow h_i = h'_i = 0$
- 2 $g_i = 1 \Rightarrow h_i = h'_i = 1$
- 3 $g_i = 2 \Rightarrow h_i = 0, h'_i = 1$ or $h_i = 1, h'_i = 0$

Example: Haplotypes $h = (1, 0, 1, 1)$ and $h' = (0, 0, 1, 0)$ explain genotype $(2, 0, 1, 2)$.

HIPP

- **Input:** Set of genotypes $G = \{g_1, g_2, \dots, g_n\}$ with $g_i \in \{0, 1, 2\}^m$.
- **Output:** Set of haplotypes $H = \{h_1, h_2, \dots, h_p\}$ with $h_j \in \{0, 1\}^m$, where
 - 1 Each genotype in G can be explained by two haplotypes in H .
 - 2 $|H| = p$ is minimal.
- **Remark:** HIPP is **NP-hard**.

Example

- **Input:** $g_1 = (2, 1)$, $g_2 = (2, 2)$.
- **Solution:** $h_1 = (0, 1)$, $h_2 = (0, 0)$, $h_3 = (1, 1)$, because
 - 1 g_1 can be explained by h_1, h_3 .
 - 2 g_2 can be explained by h_2, h_3 .
 - 3 No set of haplotypes exists with this condition and cardinality 1 or 2.

Disadvantage of HIPP

- HIPP may have **multiple** optimal solutions.
- Arbitrary algorithm returns arbitrary optimal solution.
- Found optimal solution may be **not** the **biological true solution**.

Idea

- Find **all** optimal HIPP solutions.
- **Purpose**: Increase probability of finding the biologically true solution.
- ↳ **Complete Haplotype Inference by Pure Parsimony (CHIPP)**
- ↳ Biological Study:
[Climer, J., Templeton, Zhang; Bioinformatics; 2009]

CHIPP

- **Input:** Set of genotypes $G = \{g_1, g_2, \dots, g_n\}$ with $g_i \in \{0, 1\}^m$
- **Output:** Find **all HIPP** solutions.

Example

- **Input:** $g_1 = (2, 2)$.
- **Solution 1:** $h_1 = (0, 0), h_2 = (1, 1)$.
- **Solution 2:** $h_3 = (0, 1), h_4 = (1, 0)$.

Integer Programming (IP) Model for HIPP

- [Gusfield; 2003]
- Consider **all** haplotypes h_1, h_2, \dots, h_r explaining at least one genotype.
- Define for $i = 1, 2, \dots, r$:

$$x_i = \begin{cases} 1, & \text{if haplotype } h_i \text{ appears in the HIPP solution} \\ 0, & \text{otherwise} \end{cases}$$

- Structure of IP model

$$\min \sum_{i=1}^r x_i \quad \text{subject to}$$

$$x_i \in \{0, 1\} \quad \text{for } i = 1, \dots, r,$$

further constraints

IP Algorithm for CHIPP

- 1 Solve HIPP using the IP model.

Let i_1, \dots, i_p be the indices of these haplotypes.

- 2 Add the following inequality to the IP model:

$$\sum_{s=1}^p x_{i_s} \leq p - 1$$

- 3 Solve the newly expanded IP.

- ④ **Case 1:** The new IP has an objective value larger than p .
↳ No new optimal solution exists.
- ⑤ **Case 2:** The new IP has an objective value equal to p .
↳ Another optimal solution has been found.
- ⑥ Repeat this process, until all optimal solutions have been found.

Branch-and-Bound (BnB) Algorithm for HIPP

- [Wang, Xu; 2003]
- The algorithm starts with a heuristic solution leading to the initial upper bound for the BnB search.
- The search implicitly considers all possible explaining haplotype pairs for each genotype.
- The best solution found is the optimal solution to be returned.
- If during the search the node cost is equal to or exceeds the upper bound, move on to the next branch.

BnB Algorithm for CHIPP

- Pruning is applied only when the node cost **strictly** exceeds the upper bound.
- This allows to explore a branch that may lead to **another** optimal solution.

Concept

- **Idea:** Transform instance to a smaller and easier equivalent one.
- **Advantage:** New instance easier to solve.
- **Disadvantage:** Transformation costs extra time.
- **Expectation:** For difficult instances extra time is much smaller than saved time.

↪ Fixed Parameter Tractability [Niedermeier; 2006]

Backbones

- **Backbone haplotypes:** Haplotypes appearing in each optimal HIPP solution.
 - **Backbone genotypes:** Genotypes that can be explained by two backbone haplotypes.
 - **Idea:** Backbone genotypes can be omitted for HIPP/CHIPP.
Reason: They can be explained by haplotypes of **any** optimal solution.
- ↪ Compute backbone haplotypes.
- ↪ With backbone haplotypes compute backbone genotypes (easy).

Computation of trivial backbone haplotypes

- A genotype with **no 2** leads to one backbone haplotype.
Example: $g = (1, 0, 1) \Rightarrow h = (1, 0, 1)$.
- A genotype with **one 2** leads to two backbone haplotypes.
Example: $g = (1, 0, 2) \Rightarrow h = (1, 0, 0)$ and $h' = (1, 0, 1)$.
- This are **trivial backbone haplotypes**.

Computation of non-trivial backbone haplotypes

- **Idea:** When a backbone haplotype is omitted, no optimal solution to HIPP can be found any more.
- This haplotype **must** be contained in the optimal solution.

Example

- **Input:** $g_1 = (1, 0, 2, 2)$, $g_2 = (2, 2, 1, 0)$, $g_3 = (2, 2, 1, 2)$.
- **Solution 1:** $h_1 = (1, 0, 1, 0)$, $h_2 = (1, 0, 0, 1)$,
 $h_3 = (0, 1, 1, 0)$, $h_4 = (0, 1, 1, 1)$, because
 - 1 g_1 can be explained by h_1, h_2 .
 - 2 g_2 can be explained by h_1, h_3 .
 - 3 g_3 can be explained by h_1, h_4 .
 - 4 No set of haplotypes exists with this condition and cardinality 1, 2 or 3.
- **Question:** Is $h_1 = (1, 0, 1, 0)$ a backbone haplotype?
- **Answer:** Yes, because a solution **without** $h_1 = (1, 0, 1, 0)$ contains
 $h_5 = (1, 0, 0, 0)$, $h_6 = (1, 0, 1, 1)$,
 $h_7 = (0, 0, 1, 0)$, $h_8 = (1, 1, 1, 0)$ and another 5-th haplotype.

Example

- **Input 1:** $g_1 = (2)$.

Solution: $h_1 = (0), h_2 = (1)$.

- **Input 2:** $g_2 = (2, 2)$.

Solution 1: $h_3 = (0, 0), h_4 = (1, 1)$.

Solution 2: $h_5 = (0, 1), h_6 = (1, 0)$.

Equal Column Technique

- Copy column 1 of the solution of input 1 to column 2.
- ↳ Solution 1 of input 2.
- ↳ Equal column technique for **HIPP** [Wang, Xu; 2003].
- **But:** original equal column technique does **not work** directly for **CHIPP** because solution 2 of input 2 is missed by this method.
- ↳ **Idea:** For an equal column enumerate all possible combinations of **0** and **1** with the restriction that
 - each resulting haplotype explains at least one genotype.

Decomposability Algorithm

- Two genotypes g and g' are **non-overlapping**, if $i \in \{1, 2, \dots, m\}$ exists with:
 - $g_i = 0, g'_i = 1$ or $g_i = 1, g'_i = 0$.
 - **Observation:** Two **non-overlapping genotypes** do not share any **explaining haplotype**.
- ↪ **Idea:** Compose CHIPP solutions of sub-classes whose genotypes do not overlap.

Example

- **Input 1:** $g_1 = (0, 2, 2)$.
Solution 1: $h_1 = (0, 0, 0), h_2 = (0, 1, 1)$.
Solution 2: $h_3 = (0, 0, 1), h_4 = (0, 1, 0)$.
- **Input 2:** $g_2 = (1, 2, 2)$.
Solution 1: $h_5 = (1, 0, 0), h_6 = (1, 1, 1)$.
Solution 2: $h_7 = (1, 0, 1), h_8 = (1, 1, 0)$.
- **Input 3:** $g_1 = (0, 2, 2), g_2 = (1, 2, 2)$.
Solution 1: h_1, h_2, h_5, h_6 .
Solution 2: h_1, h_2, h_7, h_8 .
Solution 3: h_3, h_4, h_5, h_6 .
Solution 4: h_3, h_4, h_7, h_8 .

Easy Instances

- **Optimized** slightly worse than **baseline** algorithms.

↪ **Typical example**: IP-Bas.: 1 sec. IP-Opt.: 1.68 sec.

Difficult Instances

- **Optimized** significantly superior to **baseline** algorithms.

↪ **Typical example**: BnB-Bas.: > 6 h. BnB-Opt.: 275.69 sec.

Optimized Versions

IP better performance than **BnB**.

Optimization Techniques

- Strongly depends on **structure** of instances, e.g., backbone technique effective for instances with many backbones.

Thanks for your attention!